

The Pynchon Gate

A Secure Method of Pseudonymous Mail Retrieval

Len Sassaman¹, Bram Cohen², and Nick Mathewson³

¹ K. U. Leuven ESAT-COSIC
Kasteelpark Arenberg 10,
B-3001 Leuven-Heverlee, Belgium
`len.sassaman@esat.kuleuven.ac.be`

² BitTorrent
227 Bellevue Way NE, Suite 152,
Bellevue, WA 98004 USA
`bram@bitconjurer.org`

³ The Free Haven Project
`nickm@freehaven.net`

Abstract. We argue for the Pynchon Gate, a practical pseudonymous message retrieval system. Our design uses a simple distributed-trust private information retrieval protocol to prevent adversaries from linking recipients to their pseudonyms, even when some of the infrastructure has been compromised. This approach resists global traffic analysis significantly better than existing deployed pseudonymous email solutions, at the cost of additional bandwidth. We examine security concerns raised by our model, and propose solutions.

1 Introduction

Pseudonymous messaging services seek to provide users with a way to send messages that originate at a pseudonymous address (or “nym”) unlinked to the user, and to receive messages sent to that address, without allowing an attacker to deduce which users are associated with which pseudonyms. These systems can be used specifically to provide a mechanism for a user to communicate without revealing her identity, or can be used as a building-block for other systems that need a bi-directional anonymous communication channel, such as Free Haven [27]. But, as we will argue below, most existing deployed solutions are either vulnerable to traffic analysis or require unacceptably large amounts of bandwidth and storage as the number of users and volume of traffic increase.

We propose the Pynchon Gate, a design that uses private information retrieval (PIR) [13] primitives to build a secure, fault-tolerant pseudonymous mail retrieval system.

In our system, pseudonymous users (or “nym holders”) use an existing anonymous email network (such as Mixmaster [48] or Mixminion [22]) to send authenticated requests to a *nym server*, which delivers outgoing messages to the email

network and handles administrative commands. The nym server also receives incoming messages and passes them to a *collator* component, which encrypts the messages and periodically packages them into regular batches. These batches are then replicated at a number of *distributor* servers, which use a private information retrieval protocol to allow nym owners to receive mail while maintaining unlinkability between a message and its recipient.

Goals. First, our design must be *secure*: we want the Pynchon Gate to resist active and passive attacks at least as well as the state of the art for forward message anonymity. Thus, we should protect users’ identities from a global eavesdropper for as long as possible; should hinder active attackers who can delay, delete, or introduce traffic; and should resist an attacker who has compromised some (but not all) of the servers on the network.

In order to provide security, however, we must ensure that the system is *deployable* and *usable*—since anonymity and pseudonymity systems hide users among each other, fewer users means less protection [1]. This implies that we should handle node failure without loss of mail; that we must not require more bandwidth than volunteer servers can provide or users are willing to use; and that we should not require a complicated interface.

In this paper. We begin in Section 2 with a discussion of related work, and an overview of known attacks against existing pseudonymity systems. (To motivate our work, Subsection 4.1 presents new analysis on the effectiveness of passive traffic analysis against current reply-block based nym servers.) Section 3 presents the Pynchon Gate in more detail, describing its organization, design rationales, and network formats. We describe our simple PIR protocol in subsection 3.4. In Section 4 we analyze security, and in Section 5 we discuss optimizations and compare our performance to that of other pseudonymous message systems. We close with an evaluation of our design in Section 6.

2 Background and attacks

Here we present a brief outline of existing pseudonymity solutions, discuss attacks against them, and present novel analysis of the effectiveness of one kind of traffic analysis against the most popular currently deployed design.

2.1 Related Work

Below we discuss existing designs for pseudonymous message delivery. Many assume the existence of a “forward” anonymous channel that a sender can use to send a message to a known recipient while preventing the recipient, the infrastructure, and any attackers from knowing who is communicating with whom. Currently deployed designs are based on Chaum’s mix [12] architecture, and include the Mixmaster [48] and Mixminion [22] anonymous remailer networks.¹

¹ Other descriptions of the use of PIR in preserving recipient anonymity have been independently proposed but not deployed [6, 15]. Independent work by Jim Mc-

It is trivial to use these systems to *send* pseudonymous messages: the sender can make an anonymous message pseudonymous by signing it with a public key associated with her pseudonym. Thus, these designs focus on how to *receive* messages sent to a pseudonymous address.

Reply blocks and return addresses. In 1981, Chaum [12] described a method of using *return addresses* in mix-nets: recipients encode a reply path, and allow senders to affix messages to the encoded path. As the message moves through the network, the path is decoded and the message encoded at each hop, until an encoded message reaches its eventual recipient. This system relies upon all selected component nodes of the chosen path remaining operational in order for mail to be delivered, which can make the system too unreliable for practical use if significant time elapses between path generation and message origination.²

In addition to reliability issues, some implementations of these “reply blocks” suffer from a pseudonym management perspective. Cypherpunk nym servers based on the first generation implementation of Chaum’s mix-nets (Type I remailers [29]), such as `alpha.c2.net` [3] and `nym.alias.net` [45], implement a central reply-block repository that allowed pseudonym holders to receive messages delivered to a email address. Unfortunately, Type I remailers allow multiple uses of their reply blocks, which are vulnerable to replay and flooding attacks as discussed in [16, 44]. Type II (Mixmaster) and Type III (Mixminion [22]) systems do not permit multiple-use reply blocks, and prevent replay attacks [17].

Single-use reply blocks. While the Type II system does not support anonymous reply blocks, the Type III (Mixminion) system introduces single-use reply blocks (SURBs) [39] to avoid replay attacks. The Type III protocol requires the recipient to create a large number of reply blocks for senders to use. In practice, this is likely to be automated by a nym server [40] that stores a number of SURBs and uses them to deliver pseudonymous mail to the recipient—one such design is Underhill [42]. Type III also has the property that the forward and reply messages share the same anonymity set, and recent work has been done by Danezis and Laurie on attack-resistant anonymous packet formats suitable for reply messages [23]. However, since reply blocks are still being used, reliability issues remain. (If any given node in the pre-selected SURB is defunct at the time mail is set to be delivered, the mail will be lost.) Reply block systems are also susceptible to intersection attacks [8]: A global observer can collect data on who is sending and receiving mail, and given enough time and data, can reliably determine who is talking to whom [19].

Network-level client anonymity. The ZKS Freedom Network [9] provided anonymous access to a POP3 server [47], enabling its users to maintain pseudonyms

Coy describes a similar architecture to the Pynchon Gate, but does not use an information-theoretic primitive for preserving privacy [46].

² Forward-only messages through a mix-net, however, are sufficiently reliable. The client software can evaluate network health information [49] before sending a message, and thus can construct robust remailer chains based on the current health of the remailer network.

using standard email protocols. Freedom was discontinued due to high operating expenses, especially in bandwidth. Other network-level anonymity systems, such as Popenet [18], Onion Routing [33], the Java Anon Proxy [7], or Tor [28], could be used in much the same fashion; unfortunately, they also suffer the same barriers to deployment [32]. Low-latency anonymity systems such as these are also far more susceptible to traffic analysis methods such as end-to-end timing attacks than are high-latency mix networks [21, 20, 25].

Network-level server anonymity. The second generation implementation of Onion Routing, Tor [28], implements rendezvous points [31] that allow users to offer location-hidden services. A user wishing to anonymously receive messages could use this to receive mail at a hidden location: Messages would be delivered to the server over the Onion Routing network, and successful delivery would not require the sender to know the IP address of the destination server.

Rendezvous points offer an alternative method of leveraging network-level anonymity systems for anonymous mail receipt; however, they do not address the previously mentioned concerns with these anonymity systems.

Re-encryption mixes. Re-encryption mixes [35] aim to improve the reliability of anonymous message systems. Recent work has shown that re-encryption mixes can be used to facilitate anonymous message replies [34]. While reusable anonymous return channels in re-encryption mixes do improve on the robustness of simple reply blocks in a Chaumian mix-net, reliability problems are still possible. Re-encryption mixes require that the security vs. reliability tradeoffs be made by the sender at the time that the message is sent. A more desirable property would be to allow the recipient to make security determinations at the time the message is retrieved.

Broadcast messages and dead-drops. Chaum discusses a traffic-analysis prevention method wherein all reply mail in the anonymous mail system is sent to all possible recipients. A less invasive optimization has already been implemented in the form of Usenet mail drops [10]: an anonymous remailer can deliver mail to a newsgroup, rather than to an email recipient. Such mail can be encrypted to a recipient’s private key, and left for her to collect from the newsgroup. If recipients use the same newsgroup and behave identically (for instance, by downloading the entire set of newsgroup messages daily), the possible statistical attacks on direct mail delivery of reply messages to individual email addresses are avoided. This solution also removes the necessity for reply-blocks, as the drop location can be established upon out-of-band.

Of course, this “send everything everywhere” approach suffers massive scalability problems. As the number of users in the system increases, each user’s bandwidth requirements become prohibitive. Users are thus encouraged to “cheat” and only download sections of the newsgroup that they are sure contain their messages, or not download on days that they do not expect messages. This allows an attacker to gather information about messages in which an individual has interest, and provides a way to attack the security of the system [2].

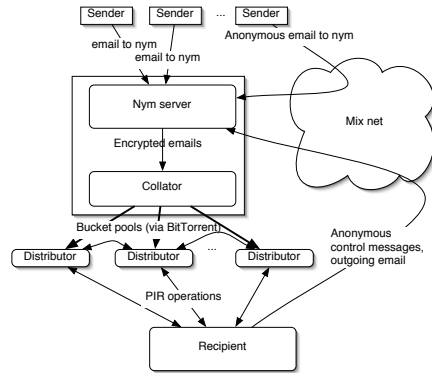


Fig. 1. The Pynchon Gate Architecture

3 The Pynchon Gate Design

We present a design framework for the Pynchon Gate. A detailed implementation specification can be found in [41].

3.1 Overview and Rationale

The Pynchon Gate is a group of servers that provide anonymous message retrieval capabilities (see Figure 1). A nym server receives messages for different pseudonym accounts via email.³ Once every “cycle” (e.g., 24 hours), the nym server passes these messages to a collator, which batches them into an indexed “bucket pool,” and passes these pools to each independently operated distributor node in the network. Each pseudonym holder then makes a series of requests to k chosen distributor nodes, enabling her to receive her messages without the individual distributors determining the pseudonym being requested. The protocol used is resistant to collusion: even if the adversary controls $(k - 1)$ of the chosen distributors, the adversary cannot link the user to her pseudonyms.

This distributed-trust PIR-based message retrieval system lets us keep the reliability, and security of the “send everything everywhere” method, while bringing the system into the realm of feasibility for users with limited resources.

We discuss the components of the Pynchon Gate architecture below.

3.2 The Nym Server

The public-facing side of the Pynchon Gate consists of a nym server that sends and receives pseudonymous email. The nym server itself provides no sender

³ The servers could also receive messages through any suitable medium for message transfer, such as “instant message” systems [24]. We require a forward anonymity protocol to allow the nym holder to communicate with the nym server, so at a minimum the nym server must be able to receive email in addition to any optional support for other protocols.

anonymity; rather, it relies on existing mix networks [22, 48]. The nym server is visible to external email correspondents, and receives messages for the nym owners at their specified email addresses.

Nym servers manage email accounts for pseudonyms. For each pseudonym, the nym server stores a *long-term public key* used by the nym holder to encrypt and authenticate outgoing email and administrative messages. Similarly, nym server stores a *short-term shared secret* for each account, used to encrypt messages to the nym holder. This secret can be reset by the nym holder after account creation.

The shared secret is updated every cycle, such that, if $S[i]$ is the shared secret in cycle i , then $S[i + 1] = H(S[i] \parallel \text{"NEXT CYCLE"})$, where $H(\cdot)$ is a cryptographic hash and \parallel denotes concatenation. From each $S[i]$, the nymserver derives a set of sub-secrets for individual messages received that cycle. The j 'th sub-secret on day i is $\text{Subkey}(j + 1, i) = H(\text{Subkey}(j, i) \parallel \text{"NEXT SECRET"})$, with $\text{Subkey}(0, i) = H(S[i] \parallel \text{"NEXT SECRET"})$.

Once it no longer needs a shared secret or a given subkey, the nym server drops it immediately, to limit the impact of key compromise (at the server or client) and improve forward security. We use a separate chain of keys for each cycle so that it is easier for a user to resynchronize after missing a few cycles.

These subkeys are used to encrypt and identify the messages received on day i . When the j 'th message for the nym is received, the nym server compresses it, encrypts it with the symmetric key $H(\text{Subkey}(j, i) \parallel \text{"KEY"})$, and prefixes it with the opaque identifier $H(\text{Subkey}(j, i) \parallel \text{"ID"})$. By deriving the message keys and identifiers in this way, we allow users to store keys and identifiers for pending messages without risking exposure of messages encrypted in the same cycle.

Finally, the nymserver also generates an different independent identifier for each user every cycle: $\text{UserID}[i] = H(S[i] \parallel \text{"USER ID"})$.

3.3 The Collator

At the end of each cycle, the nym server passes messages to the *collator*, which typically resides on the same physical server. The collator organizes all previously unretrieved messages into a three-level structure, consisting of a *meta-index*, a set of fixed-size *index buckets*, and a set of fixed-size *message buckets*.

Each user's messages are stored in a different set of message buckets, ordered by UserID. The index buckets contain, for each UserID (in order), the first message bucket containing that user's messages, and a digest of that bucket. Finally, the meta-index lists, for each index bucket, the first and last UserID in that bucket, and a digest of that bucket (see Figure 2). The index buckets and the message buckets together comprise the cycle's "bucket pool." To ensure integrity, each bucket contains a hash of the next.

The metaindex is signed with the collator's private key, along with the index of the cycle to which it applies.

To prevent an attacker from flooding a nym and observing which user receives a large volume of traffic, each nym has a maximum number of message buckets that may be filled in a given cycle. If there are more pending messages than

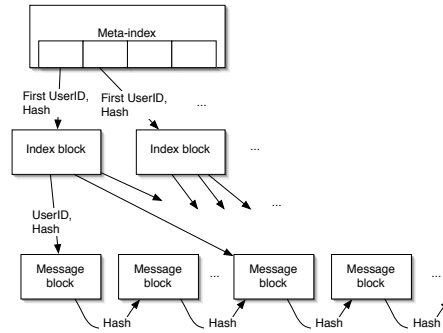


Fig. 2. The meta-index and bucket pool

will fit in the nym’s buckets, the collator defers some for a later cycle.⁴ Because the encrypted messages are prefixed with $H(\text{Subkey}(j, i) \parallel \text{ID})$, the user can tell which key to use for messages that are delivered out of order.

3.4 Distributors and clients

Once the collator is done, it relays the signed meta-index, and the entire bucket pool, to a set of independently operated *distributor nodes*. (This data should be transmitted using a bandwidth-sparing protocol such as BitTorrent [14], so that the collator does not need to send the entire pool to each distributor.)

At this point, clients can download their messages for the cycle. First, a given client downloads the meta-index from a randomly chosen distributor, and verifies its signature. The client then computes its UserID for the day, and uses the meta-index to tell which index bucket will contain an entry for that UserID.

The client then uses the PIR protocol described below to retrieve the correct index bucket, checks that the bucket’s digest is as expected, and uses the index bucket to learn which message buckets will contain the client’s messages. The client downloads these buckets with PIR, and checks their digests. If the client has received fewer buckets this cycle than her maximum, she performs extra PIR operations up to that maximum, to prevent an observer from learning how many messages she has received.

Depending on the length of the cycle, clients may not be able to download messages every cycle. Therefore, distributors must retain meta-indexes and bucket pools for a reasonable window of time, to be sure that all clients have time to download their messages.

⁴ As an extension to save bandwidth and prevent denial of service attacks, the nym server can build a special “summary message” containing the headers of pending emails, and their opaque identifiers, and include this message in the user’s message bucket. The user can then send a signed control message to the nym server requesting that unwanted emails be deleted and desired ones given priority.

Since it is not necessary for every distributor to be operational at the given point that a client wishes to retrieve mail, the system handles distributor node failure in a graceful manner.

The PIR Protocol This simple PIR protocol allows a client to retrieve a bucket from k chosen distributors, so that an attacker cannot tell which bucket the client is retrieving without compromising or controlling all k of the servers.

The protocol runs as follows: after choosing distributors, the client establishes an encrypted connection to each (e.g., using TLS [26]). These connections must be unidirectionally authenticated to prevent man-in-the-middle attacks, and can be made sequentially or in parallel.

The client sends a different “random-looking” bit vector v_{sb} to each distributor s , for each bucket b to be retrieved. Each bit vector has a length equal to the number of buckets in the pool. Each distributor s then computes $R(v_{sb})$ as the XOR of all buckets whose positions is set to 1 in v_{sb} . The resulting value is then returned to the client.

Thus, in order to retrieve the b 'th bucket, the client need only to choose the values of v_{sb} so that their exclusive OR is 0 at every position except b . (For security, $k - 1$ of the vectors should be generated randomly.) When the client receives the corresponding $R(v_{sb})$ values, she can XOR them to compute the bucket's contents.

As an optimization, a client may send $k - 1$ of the distributors a key for a stream cipher instead of a bit vector. The distributors can use the stream in place of the vector [4, 6]; only one still needs to receive a full vector.

4 Known attacks against pseudonymity systems

Most attacks on pseudonymity systems fall into one of the following categories.

Legal and hacking attacks. Attackers may attempt to coerce the operators of pseudonymity systems through lawsuits or other means [45, 52, 37, 36, 30], or may attempt to surreptitiously obtain information about nym holders.

We limit these effects of these attacks by greedily encrypting incoming messages and deleting encryption keys. All sensitive data is deleted as the bucket pool is generated, ensuring that the collator has as little information useful to an attacker as possible.

Without dynamic key rotation it would be trivial for an attacker to archive all data sent to distributors, and then at some later time obtain the nym's collator address and key from the nym server through an active attack on those components. The attacker could then read all messages that nym has ever received. In the interest of retaining little information for an attacker, implementations should discard old secrets *as soon as they are no longer needed*. Thus, at the start of each cycle i , a nymserver should derive $S[i + 1]$, $\text{UserID}[i]$, and $\text{Subkey}(0, i)$, and immediately discard $S[i]$. After using each $\text{Subkey}(j, i)$, the nymserver should calculate $\text{Subkey}(j + 1, i)$ and discard $\text{Subkey}(j, i)$. After each cycle, the nymserver should discard the last $\text{Subkey}(j, i)$, and $\text{UserID}[i]$.

Mix attacks. Since we rely on mix networks, we must be concerned with attacks against them. Furthermore, reply-block-based nym server systems require additional security properties that normal mix-net systems may not have [23].

The Pynchon Gate uses mix-nets for forward message delivery only. Attacks that do not work against a mix-net in normal forward-delivery mode will not impact the Pynchon Gate.

Man-in-the-middle attacks. An attacker able to pose as a user’s chosen distributors could trivially see all k PIR requests. We use TLS authentication to prevent this attack.

Tagging and known-cleartext attacks. An attacker may alter a message, or observe the cleartext of a message, so that he may be able to later link an input message with a given output retrieved by the nym holder.

The Pynchon Gate’s message and link encryption prevents an attacker from observing the cleartext of a message. Tagging attacks are ineffective, as TLS protects data integrity on the wire. The metaindex provides the client with the hash of the index bucket. Each message bucket provides the hash of the next message bucket, and with this information, the client can verify the integrity of information downloaded from distributors and respond to garbled data without leaking information about which bucket it was requesting.⁵

“Who am I?” attack. An attacker may send messages intended for nym Alice to nym Bob instead, so that if “Alice” responds, the attacker will know they are the same nym holder [20].

This attack relies primarily upon the ability to link one nym, $Alice_1$, with another nym, $Alice_2$, by sending a message encrypted to $Alice_1$ ’s to $Alice_2$ ’s address. The Pynchon Gate avoids this, though a similar social engineering attack may be performed if the nym holder is using a separate message encryption protocol such as PGP [11]. More research needs to be done to improve the area of privacy-preserving human-computer interaction [50, 53].

Usage pattern and intersection attacks. An attacker may analyze network usage and anonymity set members over time to sub-divide anonymity sets such that a given user is identified. In addition to passive observation of the network, there are a number of active attacks. For example, an attacker could flood a nym to observe a corresponding increase in traffic by the recipient.

Users of the Pynchon Gate select distributors from which to receive data at random, each time the nym holder retrieves her messages. Unlike systems where the pseudonym infrastructure initiates the delivery of messages, the Pynchon Gate Client initiates the retrieval of messages, and thus message retrieval cannot be correlated to a given nym by a malicious sender.

Message buckets are of a fixed size, to protect against active flooding attacks [51] as well as simple usage pattern analysis. If the volume of messages is

⁵ If a client *does* notice a corrupt bucket, it should not re-attempt the PIR operation until it has received all buckets, to avoid leaking which bucket it was reading through the timing of its response.

too great to fit in a users’ buckets, delivery continues by *trickling* the pending messages to the distributors over the next several tree distributions. To prevent denial of service attacks, users can selectively retrieve or delete excess messages.⁶

Since the time between sending a message and receiving a reply may leak information about the nym holder, traffic from the client to the distributors is regulated by the client, which queries the distributors only at given intervals. To thwart active attacks against the distributor targetting a specific client, clients should choose these intervals randomly.

4.1 Statistical disclosure against reply-block-based nym servers

Nym servers based on reply blocks (discussed in Section 2.1 above) are currently the most popular option for receiving messages pseudonymously. Nevertheless, they are especially vulnerable to end-to-end traffic analysis.

Suppose an adversary is eavesdropping on the nym server, and on all recipients. The attacker wants to know which user (call her Alice) is associated with a given pseudonym (say, `nym33`). The adversary can mount an *intersection attack*, by noticing that Alice receives more messages, on average, after the nym server has received a message for `nym33` than when it has not.⁷ Over time, the adversary will notice that this correlation holds for Alice but not for other users, and deduce that Alice is likely associated with `nym33`.

Recent work [19, 43] has studied an implementation of these intersection attacks called *statistical disclosure*, where an attacker compares network behavior when Alice has sent to network when she is absent in order to link an *anonymous* sender Alice to her regular recipients `Bob1...BobN`. Against *pseudonymous* recipients, however, these attacks are far easier: in the anonymity case, many senders may send to any given recipient `Bobi`, but with pseudonymous delivery, only one user sends or receives messages for a given pseudonym.

To examine this effect, we ran a version of the attack simulations described in [43], assuming a single target pseudonym and N non-target pseudonyms providing cover. In order to make the attack as difficult as possible (and thus establish an upper bound for security), we assume that users behave identically: they receive messages with equal probability according to independent geometric distributions in each unit of time (receiving no messages with probability $1 - P_M$); they use identical reply blocks with path length ℓ through mixes in a steady state that delay each message each round with probability P_D .

We ran the simulated attack with different values for P_M , P_D , and ℓ , against a nym server with $N = 2^{16}$ active pseudonymous users. (This is probably an overestimate of the number of users on typical nymserver today [45].) We performed

⁶ If a client will be retrieving large amounts of data on a regular basis, this method will not work. Instead, the client should at account creation time request a sufficient number of buckets to receive all data destined to it. Pending data queued on the collator should be expired after a reasonable amount of time.

⁷ This task is especially easy if the adversary can distinguish reply messages from non-reply messages, as is possible with Type I remailers.

100 trials for each set of parameters. In the worst case (for the nym holder), when $P_M = 0.5, \ell = 1, P_D = 0.1$, the lack of mix-net delay variance allowed the simulated attacker to guess the user’s identity correctly after the user received an average of only 37 messages. In the best simulated case ($P_M = 0.5, P_D = 0.9, \ell = 4$), the user received an average of only 1775 messages before the attacker guessed correctly. For an active user, this is well within a month’s expected traffic.

Although there are ways to use dummy traffic to resist statistical disclosure attacks, these difficult to implement perfectly in practice (due to outages) and even slight imperfections render users vulnerable to attack [43].

4.2 Other security concerns

The information used for authentication of the system components (such as the certificates and the hash tree root and metaindex) must be published widely to prevent either the collator or any of the distributors from attacking clients by tricking them into thinking that the hash root of the tree is something other than what all of the other clients know it to be. Distributors should do basic sanity checks, such as verifying tree integrity. The distributors should also send audit messages of their own to verify that the messages correctly appear in the system. Finally, clients should make sure that each of the distributors they use agree about the value of the hash root.

5 Performance, Scalability and Optimizations

In this protocol, the size of requests is linearly proportional to the total number of messages and the size of responses is the bucket size. If one or the other of these values is large enough to cause scaling problems, then the collator can trade off bucket size for bit field size by doubling the bucket size, which halves the bit field size. With this approach, if the number of buckets becomes very large, then the message size rises proportionately to the square root of the number of buckets. This scales well, although it may necessitate multiple collators if the number of buckets gets extremely large. (Note that while different collators may share the same distribution nodes for architecture or resource reasons, their anonymity sets would remain separate.)

The PIR algorithm suggested in this paper does not have optimal asymptotic performance, especially in bandwidth. We present it nonetheless because it is easy to explain, implement, and analyze, and offers sufficient scalability to be useful. A reasonable engineering plan is to implement the algorithm that we describe, then once the implementation reaches a level of scaling high enough to warrant the much greater difficulty of a more sophisticated algorithm, implement a follow-on protocol that uses fewer resources. This delay is prudent, since private information retrieval primitives are an area of active research with ongoing improvements [5], so waiting to implement a more sophisticated algorithm will likely result in greater resource savings once the implementation occurs.

System	Nymserver bandwidth	Infrastructure bandwidth	User bandwidth	Nymserver storage
Type I nymserver	$\sum \text{Vol}_i + \text{CVol}_i$	CVol_i	$\frac{2\ell \sum \text{CVol}_i}{S}$	rN
Type III nymserver ^a	$\sum \text{Vol}_i + (M+r) \sum \left\lceil \frac{\text{CVol}_i}{P} \right\rceil$	$\frac{2L(M+r)}{S} \sum \left\lceil \frac{\text{CVol}_i}{P} \right\rceil$	$(P+r) \left\lceil \frac{\text{CVol}_i}{P} \right\rceil$	$rW \sum \left\lceil \frac{\text{CVol}_i}{P} \right\rceil$ (best case)
Usenet drop	n/a	$\frac{W}{S} \sum \text{CVol}_i$	$\left\lceil \frac{N}{S} + 1 \right\rceil \sum \text{CVol}_i$	$\sum \text{CVol}_i$ n/a
The Pynchon Gate	$\sum \text{Vol}_i + \text{Pool}$	$\frac{1}{8} [\sum \text{ClientB}_i + \text{Pool}]$	$2MEI + \text{ClientPIRVol}^b$	$W\text{Pool}$

^a Underhill can be used in a full padding mode. In this case, the performance evaluation is the same, except that CVol_i is calculated as the maximum compressed volume a user can receive, rather than the average.

^b ClientPIRVol is the amount of data sent and received during PIR, or $\text{Buckets}_i \left[(K-1)SS + \frac{(m+I)}{8} + B \right]$

Fig. 3. Performance comparison for several pseudonymity designs.

Another potential bottleneck lies in the fact that distributors have to perform a linear scan of the entire bucket pool in order to fulfill a request. However, they can use a single linear pass to compute the results of many requests in parallel. Thus, a distributor can fulfill a large number of requests at once, though the latency to answer those requests is limited by the total size of the bucket pool, and the throughput to the distributor’s hard drive (unless the bucket pool fits in RAM). Also, by performing continuous linear scans of the entire database, the distributor can begin computing the result of a request at any point during the scan, finishing when the next scan returns to that same point. Thus, the latency is exactly the time of one full scan.

5.1 Comparing The Pynchon Gate to other systems

We have evaluated the resource requirements of various pseudonymity systems described in Section 2.1, and compare their respective performance in Figure 3. Bandwidth requirements for the independent components of the pseudonym system are averages per cycle. We use the term “infrastructure” to denote mix nodes in the Type I (Cypherpunk) and Type III (Underhill [42]) nym server systems, NNTP servers [38] for the Usenet news drop, and distributors in the Pynchon Gate. N is the total number of users in the system. Vol_i is the volume of messages received by user i on a given day. CVol_i is this volume after compression. S is the number of nodes in the infrastructure.

Cypherpunk nym system parameters are denoted as r for reply block size, and ℓ for the number of mixes in the reply block.

Underhill uses a payload P of size 28 KB, a reply block r of size 2 KB, and a packet size M of 32 KB. For Underhill, W is the maximum interval at which users must replenish reply blocks. Similarly, W is the window of time (in days) in which users must retrieve their mail in the Pynchon Gate.

The Pynchon Gate allocates buckets of size B . The number of message buckets needed (m) is calculated as $\sum \left\lceil \frac{\text{CVol}_i}{B} \right\rceil$ and the number of index buckets needed (I) is calculated as $\left\lceil \frac{N \cdot IE}{B - IE} \right\rceil$, where IE is the index entry size. ME is the system’s metaindex. The PIR stream seed size is SS , and K is the number of distributors chosen from which to retrieve data.

If we assume one cycle per day with 10,000 users per collator who receive 100,000 KB of data each cycle, each collator will transfer 2 GB per day. This

allows a bucket size of 10 KB, with distributor query sizes of 12.5 KB, thus allowing reasonable bandwidth requirements for both users and system operators.

6 Conclusions

We have presented a system for anonymous message retrieval that provides stronger anonymity assurance and greater robustness than other theorized or deployed high-latency pseudonymous message retrieval systems. Our system resists traffic analysis better than current deployed systems, and offers convenient scalability options.

We have proposed a client protocol that does not rely upon an obtrusive user interface. Much work remains in the field of effective user interface design for privacy and anonymity systems, particularly when the privacy component is viewed by the user as optional.

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